Some slides for 17th Lecture, Algebra 1

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16-11-2011

Public-key crytography (from Wikipedia)

- The key used to encrypt a message is not the same as the key used to decrypt it.
- Each user has a pair of cryptographic keys-a public key and a private key. The private key is kept secret, while the public key may be widely distributed.
- Messages are encrypted with the recipient's public key and can only be decrypted with the corresponding private key.
- The keys are related mathematically, but the private key cannot feasibly (ie, in actual or projected practice) be derived from the public key.
- The discovery of algorithms that could produce public/private key pairs revolutionized the practice of cryptography beginning in the middle 1970s.

ElGamal

Based on discrete logarithm problem:

Given a prime p and $y, g \in \mathbb{N}$, find x such that

$$y \equiv g^x \pmod p$$

1. Alice and Bob choose p, a big prime, and $g \in \mathbb{N}$ s.t. $0 < g < p$ and g has order $p - 1$ in $(\mathbb{Z}/p\mathbb{Z})^*$ (a generator of $(\mathbb{Z}/p\mathbb{Z})^*$)
2. Alice chooses a, with $0 < a < p$ and computes $[g^a]_p$.
   Secret Key=a
   Public Key=$[g^a]_p$
3. Bob chooses b with $0 < b < p$ and computes $[g^b]_p$.
   Secret Key=b
   Public Key=$[g^b]_p$

4. Alice wants to send a message m, $0 < m < p$ to Bob. She sends:
$$\left([g^a]_p, [m(g^b)^a]_p\right)$$
5. Bob gets $([x_1]_p, [x_2]_p)$ and computes
$$[x_2]_p([x_1^b]_p)^{-1} = [mg^{ab}]_p([g^{ab}]_p)^{-1} = [m]_p$$
and since $m < p$ he can recover m.

To encrypt the message one uses the public key of the receiver and the secret key of the sender.

6. Eve?: she had to compute b from $[g^b]_p$